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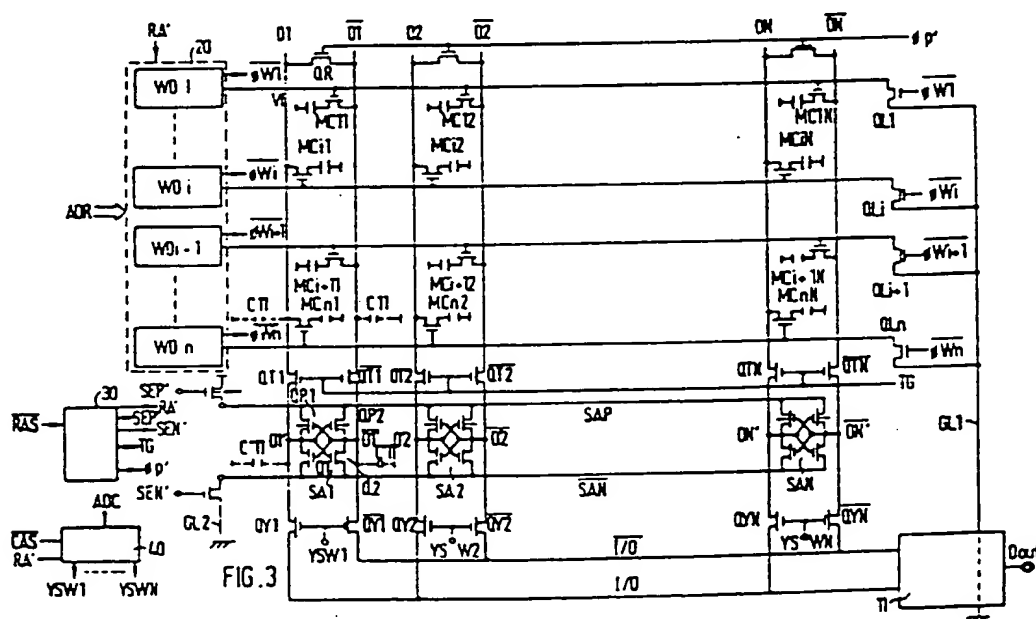
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DOG

(54) Semiconductor memory circuit having an improved restoring control circuit.

(57) A dynamic memory circuit which can operate with a reduced amount of current noise and without destruction of stored data is disclosed. The memory circuit includes dynamic memory cells necessitating restoring operation, a read circuit for performing a read-out operation in an active state and a restore circuit for performing a restore operation in a reset state following the active state.



EP 0 373 672 A2

# SEMICONDUCTOR MEMORY CIRCUIT HAVING AN IMPROVED RESTORING CONTROL CIRCUIT

## BACKGROUND OF THE INVENTION

### Field of the Invention:

The present invention relates to a semiconductor memory circuit and, more particularly to a dynamic type random access memory (DRAM) fabricated on a semiconductor substrate.

### Description of the Related Art:

Dynamic memory circuits have been widely utilized as large capacity semiconductor memories in various fields. The dynamic memory circuit is generally constructed in such a manner that one-transistor type memory cells are arranged in a matrix form of rows and columns together with word lines and bit lines arranged in rows and columns, respectively.

According to the conventional technique, the number of memory cells connected to one bit line increases as the memory capacity increases. Consequently, a stray capacitance of each bit line increases to cause various problems. First, when the sense amplifiers are activated, the bit lines need to be charged or discharged to a power supply potential  $V_{cc}$  or a ground potential  $V_{ss}$ . In this regard, the increase in the stray capacitances of the bit lines cause an increase in the time required for charging and discharging. If the respective pairs of bit lines are not sufficiently amplified, that is, if charging or discharging is not sufficiently effected, it is impossible to enable a column selection circuit so as to connect the selected pair of bit lines to a pair of common data lines. This is because, when the common data lines are connected to bit lines, the charges on the common data lines may flow into the bit lines to destroy the stored information. In consequence, the time at which the data is output from an output terminal is delayed, which results in a lower operation speed. Secondly, the increase in the bit line stray capacitances causes an increase in the charging and discharging current during an active period, thus causing problems, for example, floating of the ground potential at an internal ground wiring, a lowering in the power supply potential, and generation of noise between the bit lines. Particularly, the floating of the ground potential during the active period causes a noise due to operations of peripheral circuits such as an output circuit, the column selection circuit, and a potential of the non-selected word lines undesirably exceeds a thresh-

old voltage of memory cell transistors or more. storage capacitors of the non-selected memory cells are erroneously connected to the bit lines, resulting in destruction of the stored data.

## SUMMARY OF THE INVENTION

It is an object of the present invention to provide a semiconductor memory circuit which can operate at a high speed.

It is another object of the present invention to provide a dynamic memory circuit which can stably operate without data destruction.

The dynamic memory according to the present invention comprises means for receiving an external control signal having a first level which sets the memory in an active state and a second level which sets the memory in a reset state and dynamic memory cells which require restoring read-out data therein and, is featured in that read-out of data from a memory cell or cells is performed in the active state and restoring of data to the memory cell or cells is conducted in the reset state.

According to a preferred aspect of the present invention, a time required for the restoring of data to the memory cells is removed in the active state can be shortened and a current amount required in the active state is reduced.

Moreover, bit lines are not charged or discharged in the active state but left near a precharged level with small signal difference, memory cell transistors of the memory cells coupled to the non-selected word lines are never turned conductive even when the potential of the non-selected word lines float in levels. Thus, destruction of data stored in the non-selected memory cells can be effectively prevented.

## BRIEF DESCRIPTION OF THE DRAWINGS

The above and further objects, features and advantages of the present invention will become more apparent from the following detailed description taken in conjunction with the accompanying drawings, wherein:

Fig. 1 is a schematic circuit diagram showing a dynamic memory in the prior art;

Fig. 2 is a timing diagram showing an operation of the memory of Fig. 1;

Fig. 3 is a schematic circuit diagram of a dynamic memory according to a first embodiment of the present invention;

Fig. 4 is a timing diagram showing an opera-

tion of the memory of Fig. 3:

Fig. 5 is a schematic block diagram of a timing signal generator employed in the memory of Fig. 3:

Fig. 6 is a timing diagram showing waveforms of signals in the timing signal generator of Fig. 5:

Fig. 7 is a schematic circuit diagram of a dynamic memory according to a second embodiment of the present invention; and

Fig. 8 is a timing diagram showing an operation of the memory of Fig. 7.

## DETAILED DESCRIPTION OF THE INVENTION

### Description of the Prior Art:

With reference to Figs. 1 and 2, a typical structure of the conventional dynamic memory will be explained.

As shown in Fig. 1, the memory comprises a plurality of one-transistor type memory cells  $MC_{11}$  -  $MC_{NN}$  each having a memory cell transistor  $Q_M$  and a storage capacitor  $C_M$  connected between the transistor  $Q_M$  and a fixed potential  $V_1$ , word lines  $WL_1$  -  $WL_N$  arranged in rows, a plurality of bit line pairs  $D_1$ ,  $\bar{D}_1$  -  $D_N$ ,  $\bar{D}_N$  arranged in columns, CMOS type sense amplifiers  $SA_1$  -  $SA_N$  provided for the bit line pairs  $D_1$ ,  $\bar{D}_1$  -  $D_N$ ,  $\bar{D}_N$ , respectively, a pair of common data lines  $VO$ ,  $\bar{VO}$ , a plurality of pairs of column selection transistors  $QY_1$ ,  $QY_1'$  -  $QY_N$ ,  $QY_N'$  connected between the bit line pairs and the pair of common data lines, respectively, a row decoder 20 having decoding units  $WD_1$  -  $WD_N$  for selecting the word lines  $WL_1$  -  $WL_N$ , and an output circuit 11.

The row decoder 20 operatively selects one of the word lines  $WL_1$  -  $WL_N$  in accordance with row address signals ADR in response to a timing signal RA. Each of the sense amplifiers  $SA_1$  -  $SA_N$  includes P-channel MOS transistors  $QP_1$ ,  $QP_2$  and N-channel MOS transistors  $Q_1$ ,  $Q_2$  and is connected to a first activation line SAP and a second activation signal  $\bar{SAN}$ . The activation line SAP is connected to a p-channel MOS transistor  $QP_1$  controlled by a control signal SEP and the activation line  $\bar{SAN}$  is coupled to an N-channel MOS transistor  $Q_1$  controlled by a control signal SEN. The pairs of column selection transistors  $QY_1$ ,  $QY_1'$  -  $QY_N$ ,  $QY_N'$  are controlled by column decoder signals  $YSW_1$  -  $YSW_N$  generated by a column decoder (not shown) in a known way.

The row decoder 20 also generates clamp signals  $\phi_{w1}$  -  $\phi_{wn}$  which are opposite to those outputs for the word lines  $WL_1$  -  $WL_N$ , respectively. For example, when the word line  $WL_1$  is selected (high level), the signals  $\phi_{w2}$  -  $\phi_{wn}$  are high in level with

$\phi_{w1}$  at a low level. Clamp transistors of N-channel type  $QL_1$  -  $QL_N$  are connected between the word lines  $WL_1$  -  $WL_N$  and a ground wiring  $GL_1$  which is extending through the output circuit 11 and connected to the output circuit 11. The clamp transistors  $QL_1$  -  $QL_N$  receive the clamp signals  $\phi_{w1}$  -  $\phi_{wn}$ , respectively. While the source of the transistor  $Q_{10}$  is connected to a separate ground wiring  $GL_2$ .

With reference to Fig. 2, a typical operation of the memory of Fig. 1 will be explained.

During a reset period TP: before the starting ( $t_1$ ) of the active operation period ( $T_A$ ), the bit lines  $D_1$ ,  $\bar{D}_1$  to  $D_N$ ,  $\bar{D}_N$  have been maintained at, for example, an intermediate potential ( $1/2 V_{cc}$ ) of a power supply potential  $V_{cc}$  and a ground potential  $V_{ss}$ . When the row address strobe signal  $\bar{RAS}$  becomes an active, low level at  $t_1$ , the active period  $T_A$  is initiated and one word line driving circuit, for example,  $WD_i$ , is selected among the word line driving circuits  $WD_1$  to  $WD_N$  on the basis of row address signals ADR given externally in response to RA which is an internal signal formed front  $\bar{RAS}$  at  $t_2$ , so that one selected word line e.g.  $WL_i$  rises in potential. In consequence, the memory cell transistors  $Q_M$  of the memory cells  $MC_{i1}$  to  $MC_{iN}$  turn ON, so that the potentials of one bit lines of the bit line pairs  $D_1$ ,  $\bar{D}_1$  to  $D_N$ ,  $\bar{D}_N$  change in the range of from about 100 mV to 200 mV in accordance with charges stored in the respective storage capacitors  $C_M$ . On the other hand, the potentials of the other bit lines remain at the initial value, that is, the potential of  $1/2 V_{cc}$ . Accordingly, the signals representative of the potential differences are input to the sense amplifiers  $SA_1$  to  $SA_N$ , respectively. Next, when the sense amplifier enabling signal SEN rises and SEP falls at  $t_3$ , the sense amplifiers  $SA_1$  to  $SA_N$  are activated to amplify the very small difference signals input to the sense amplifiers. As a result, for example, the bit line  $D_1$  lowers to the ground potential, while the bit line  $\bar{D}_1$  rises to the power supply potential  $V_{cc}$  (in the case where the memory cell  $MC_{i1}$  has "0" information (L) stored therein). In actual fact, the information in the capacitor  $C_M$  of the memory cell  $MC_{i1}$  is destroyed through the exchange of charges with the bit line  $D_1$  when the word line  $WL_i$  rises to read out information to the bit line  $D_1$ . Therefore, the above-described sense amplifier operation serves not only to amplify the very small signals on the bit lines  $D_1$  and  $\bar{D}_1$  but also to rewrite (i.e., refresh) the destroyed information to the memory cell by reading out the signals to the bit lines. Thereafter, the column decoder signal  $YSW$ , which is an internal signal generated from a column address strobe  $\bar{CAS}$  on the basis of column address signals (not shown) given to an external address terminal in response to the fall of the column address strobe  $\bar{CAS}$ , rises, so that, for example,  $YSW_1$  is selected

and rises at  $t_5$ . As a result, the data from the bit line pair  $D_1$  and  $\overline{D_1}$  is transferred to the common data line pair  $I/O$  and  $\overline{I/O}$ . The transferred data is amplified by the output circuit 11 and output to the outside from the output terminal  $D_{out}$ . Then, the active period  $T_A$  is terminated at  $t_6$  and the reset signal  $\phi_p$  rises at  $t_7$ .

In the above-described prior art, as the storage capacity increases, the number of memory cells connected to each bit line increases and consequently the stray capacitances  $C_{31}$  of the bit lines increase, thus causing various problems. First, when the sense amplifiers are activated as described above, the bit lines  $D_1$  to  $D_N$  and  $\overline{D_1}$  to  $\overline{D_N}$  need to be charged or discharged to the power supply potential  $V_{cc}$  or the ground potential  $V_{ss}$ ; in this regard, the increase in the stray capacitances  $C_{31}$  of the bit lines cause an increase in the time required for charging and discharging. If the bit line pairs  $D_1$ ,  $\overline{D_1}$  to  $D_N$ ,  $\overline{D_N}$  are not sufficiently amplified, that is, if charging or discharging is not sufficiently effected, it is impossible to enable YSW so as to connect bit lines to the common data lines  $I/O$  and  $\overline{I/O}$ . This is because, when the common data lines are connected to bit lines, the charges on the common data lines may flow into the bit lines to destroy the information. In consequence, the time at which the data is output from the output terminal  $D_{out}$  is delayed, which results in the performance being deteriorated. Secondly, the increase in the bit line stray capacitances  $C_{31}$  invites an increase in the charging and discharging current, thus causing problems, for example, floating of the ground potential, a lowering in the power supply potential, and generation of noise between the bit lines.

Moreover, in the above-described conventional semiconductor memory, either the bit line  $D_1$  or  $\overline{D_1}$  is discharged to the ground potential during the active period  $T_A$  when the row address strobe signal  $\overline{RAS}$  is at the low potential. Therefore, if a noise is generated in the internal grounding wiring  $GL_1$  due to the activation of the output circuit 11 during this period, the potentials of the wiring  $GL_1$  and the non-selected word lines, for example,  $WLi + 1$ , undesirably rises from the ground potential to a potential near the threshold voltage of the memory cell transistors  $Q_M$ , while the bit line  $D_1$  already becomes the ground potential, and if the memory cell  $MCi + 11$  has information "1" (H) stored therein, the sub-threshold current of the MOS transistor in the memory cell  $MCi + 11$  causes the charge in the capacitor of the memory cell  $MCi + 11$  to flow out to the bit line  $D_1$ , thus causing destruction of the stored information. Since this phenomenon is caused by the sub-threshold current of the MOS transistor in the memory cell, failures are likely to occur frequently due to variations in manufacture, for example, misalignment,

variations in the gate length of the gate electrodes, etc. The worst is the case where the row address strobe signal  $\overline{RAS}$  is maintained at the low potential for a long period of time and during this period the column address strobe signal  $\overline{CAS}$  and the external address signal change at high frequency. Thus, a great deal of time is required for inspection.

## 10 Description of Preferred Embodiments:

With reference to Figs. 3 to 6, the dynamic memory according to one embodiment of the present invention will be explained. In Figs. 3 to 6, elements or portions corresponding to those in Figs. 1 and 2 are denoted by the same or similar references and detailed description therefor will be omitted.

The memory of this embodiment is featured as follows. Namely, a plurality of pairs of transfer gate (N-channel) transistors  $QT_1$ ,  $\overline{QT_1}$  -  $QT_N$ ,  $\overline{QT_N}$  are inserted between the pairs of bit lines  $D_1$ ,  $\overline{D_1}$  -  $D_N$ ,  $\overline{D_N}$  and a plurality of pairs of sense nodes  $D_1'$ ,  $\overline{D_1'}$  -  $D_N'$ ,  $\overline{D_N'}$  of the sense amplifiers  $SA_1$  -  $SA_N$ , respectively. The purpose of provision of the transfer gates  $QT_1$  to  $\overline{QT_N}$  is to cut off the bit line stray capacitances  $C_{11}$  from the stray capacitances  $C_{11}$  of the sense nodes  $D_1'$  to  $D_N'$  when the sense amplifiers are activated, thereby increasing the operating speed of the sense amplifiers.

A timing signal generator 30 receives the row address strobe signal  $\overline{RAS}$  and generates timing signals  $RA'$ ,  $SEP'$ ,  $SEN'$ ,  $TG$  and  $\phi_p'$ . A column decoder 40 receives column address signals  $ADC$  and generates the column decoder signals  $YSW_1$  -  $YSW_N$  under control of  $\overline{CAS}$  and  $RA'$ .

In addition, a restoring or rewriting operation is conducted during a reset period when  $\overline{RAS}$  is made inactive under control of  $RA'$ ,  $SEP'$ ,  $SEN'$ ,  $TG$  and  $\phi_p'$  generated by the timing signal generator 30.

The operation of the memory of Fig. 3 will be explained with reference to Fig. 4.

Prior to a time point  $t_1$ , a reset period  $TP_1$  has been set and the pairs of bit lines  $D_1$ ,  $\overline{D_1}$  -  $D_N$ ,  $\overline{D_N}$  and the sense nodes  $D_1'$ ,  $\overline{D_1'}$  -  $D_N'$ ,  $\overline{D_N'}$  have been precharged to a precharge potential ( $1/2 V_{cc}$  level). At  $t_1$ , the row address strobe signal  $\overline{RAS}$  becomes active (low level) to introduce the active period the row address signals  $ADR$  are taken in, and on the basis of this address signals one word line driving unit, for example,  $WD_1$ , is selected from among the word line driving circuits  $WD_1$  to  $WD_N$ . Thereafter, as the internal signal  $RA'$  formed on the basis of  $\overline{RAS}$  rises at  $t_2$ , the potential of the selected word line  $WLi$  rises and consequently information stored in the memory cells  $MCi_1$  to  $MCi_N$  connected to the selected word line  $WLi$

appear as very small signals on the bit line pairs  $D_1, \overline{D_1}$  to  $D_N, \overline{D_N}$ , respectively. Since at this time the transfer gate control signal TG is at the high potential, the signals on the bit line pairs  $D_1, \overline{D_1}$  to  $D_N, \overline{D_N}$  are also transferred to the sense nodes  $D_1, \overline{D_1}$  to  $D_N, \overline{D_N}$  respectively. Next, at  $t_2$  the control signal TG is lowered to turn OFF the transfer gates  $Q_{T1}$  to  $Q_{TN}$ , thereby cutting off the bit lines  $D_1$  to  $\overline{D_N}$  from the sense nodes  $D_1$  to  $\overline{D_N}$ . Thereafter, at  $t_3$  the sense amplifier enabling signals SEN and SEP are raised and lowered, respectively, to activate the sense amplifiers  $SA_1$  to  $SA_N$ . Since the sense nodes  $D_1$  to  $D_N$  have relatively small stray capacitances  $C_{11}$ , as described above, they are amplified at high speed. As a result, for example, one sense node  $D_1$  reaches the power supply potential  $V_{cc}$ , while the other sense node  $\overline{D_1}$  reaches the ground potential, rapidly. Thereafter, the column selection transistors, for example,  $Q_{V1}$  and  $\overline{Q_{V1}}$ , selected after the column address strobe signal  $\overline{CAS}$  becomes active at  $t_5$  on the basis of the externally applied column address information ADC, are driven to turn ON in response to the rise of the internal signal YSW at  $t_5$  generated on the basis of  $\overline{CAS}$ , so that the data from the sense nodes  $D_1$  to  $\overline{D_1}$  is transferred to the common data line pair  $I/O$  and  $\overline{I/O}$  and output from the output terminal Dout via the output circuit 11. At this point of time, the transfer gate control signal TG remains at the low potential and therefore the bit lines  $D_1$  to  $\overline{D_N}$  are not charged nor discharged but maintained near the middle potential.

Therefore, it is guaranteed that the potentials of all the bit lines are near the  $1/2 V_{cc}$  level and higher than the non-selected word lines even if the potential of the ground wiring GL<sub>1</sub> rises to some extent due to the operation of the output circuit 11. Accordingly, the memory cell transistors  $Q_M$  of the memory cells coupled to the non-selected word lines are never turned ON.

Thereafter, when  $\overline{RAS}$  and  $\overline{CAS}$  become inactive at  $t_8$  to introduce to a reset period  $TP_2$ , the internal signal YSW is first lowered to turn OFF the column selection transistors  $Q_{V1}$  and  $\overline{Q_{V1}}$ , thereby cutting off the sense nodes  $D_1$  and  $\overline{D_1}$  from the common data line pair  $I/O$  and  $\overline{I/O}$ . Thereafter, at  $t_9$  the transfer gate control signal TG is raised to start charging and discharging of the bit lines  $D_1$  to  $\overline{D_N}$  by the sense amplifiers  $SA_1$  to  $SA_N$  which have been left activated. Thus, one of each bit line pair  $D_1, \overline{D_1}$  to  $D_N, \overline{D_N}$  reaches the power supply potential  $V_{cc}$ , while the other reaches the ground potential, performing a restore operation. Thereafter, the word line driving signal RA is lowered to lower the selected word line WLi. At this point of time, the refreshing of the memory cells MCi1 to MCiN is completed. Thereafter, the sense amplifier enabling signals SEN and SEP are reset at  $t_{10}$  and the bit

line sense nodes are also reset, thus completing the operation.

In the foregoing, the arrangement in which N-type MOS transistors are employed as switching elements and CMOS type sense amplifiers are used has been described. However, various changes and modifications may be made to matters of design, for example, the selection of the above-mentioned parts and the set values for the initial and finally reached potentials, of the bit lines according to circumstances. The subject matter of the present application resides in that after the information stored in a memory cell has been transferred to the sense nodes during the active period, the sense nodes and the bit lines are left cut off from each other so as not to carry out charging or discharging of the bit lines. According to the present embodiment, the restoring data to the memory cells is conducted during the reset period after the active period  $T_A$ , and therefore, the effective length of the active period is shortened as viewed from the outside of the memory, thus, data rate in the access operation can be improved.

Moreover, the memory cell transistors  $Q_M$  of the non-selected memory cells are kept non-conductive by performing the restore operation in the reset period, destruction of stored data can be effectively prevented.

An example of the timing signal generator 30 in Fig. 3 will be explained with reference to Figs. 5 and 6.

As shown in Fig. 5, the timing signal generator 30 includes delay circuits D1 to D7, inverting delay circuits ID1 and ID2, inverter circuits IV<sub>1</sub> to IV<sub>6</sub>, a dummy row address buffer 51, NAND gates 52, 54, 55, 56 and 57, a NOR gate 53, and a series circuit of P-channel MOS transistor  $Q_{P5}$  and N-channel MOS transistors  $Q_{S1}, Q_{S2}$ .

An internal signal  $\phi_1$  is generated from  $\overline{RAS}$  through the delay circuit D1 and the inverter IV<sub>2</sub> and has the substantially opposite phase to  $\overline{RAS}$ . A control signal AE which is used to enable row address buffers (not shown in Fig. 3) is generated from the delay circuit D3. The dummy row address buffer 51 receives one of row address signals Ai and generates its true and complementary buffered signals Xi,  $\overline{Xi}$  when enabled by AE. The signals Xi,  $\overline{Xi}$  are applied to the NOR gate 53, the inverter IV<sub>3</sub> and the delay circuit D4 in cascade and, an internal signal  $\phi_2$  is generated when the states of Xi and  $\overline{Xi}$  are established. The internal signals  $\phi_1$  and  $\phi_2$  are applied to the NAND gate 54 for generation of RA and TG, as illustrated.

Internal waveforms of the respective signals are shown in Fig. 6.

Fig. 7 shows the arrangement of a second embodiment of the present invention. This arrangement differs from the arrangement shown in Fig. 3

in that bit lines are divided into two groups  $d_{1a}, \overline{d_{1a}}, d_{Na}, \overline{d_{Na}}$  and  $d_{1b}, \overline{d_{1b}}, d_{Nb}, \overline{d_{Nb}}$ , which are separated from each other through second transfer gates  $Q_{T1}, \overline{Q_{T1}}$  to  $Q_{TN}, \overline{Q_{TN}}$ .

The operation is substantially the same as that of the foregoing first embodiment except for the following. Namely, at the time when the transfer gate control signal TG is lowered at  $t_3$  in the first embodiment in the active period  $T_A$ , the first and second transfer gate control signals TG and TG2 of the second embodiment are lowered simultaneously or successively in the mentioned order and thereafter both maintained at the low potential during the active period  $T_A$  to maintain the bit lines  $d_{1a}, \overline{d_{1a}}$  to  $d_{Na}, \overline{d_{Na}}$  and  $d_{1b}, \overline{d_{1b}}$  to  $d_{Nb}, \overline{d_{Nb}}$  near the middle potential. In the reset period TP2 introduced by the fall of  $\overline{RAS}$  at  $t_4$ , the control signal TG is raised and the control signal TG2 is raised with a little delay at  $t_2$ . In other words, the feature of the second embodiment resides in that the bit lines are divided and charged or discharged in a time division manner at the time of refreshing. Thus, it is possible to reduce the generation of noise.

As has been described above, according to the present invention, after the information stored in a memory cell has been transferred to the sense nodes during the active period, the sense nodes and the bit lines are cut off from each other so as to maintain the bit line potential near the middle potential between the power supply potential and the ground potential, thereby making it possible to prevent destruction of the information stored in non-selected memory cells due to floating or sinking of the non-selected word lines caused by the ground potential line noise in the chip that is attendant on the operations of other circuit blocks on the same chip, for example, the  $\overline{CAS}$  circuit block, column address buffer and column address decoder circuit, during the active period.

Further, since the margin with respect to the sub-threshold current of the MOS transistors of the memory cells increases, the allowance for variations in manufacture is made. It is therefore possible to expect an increase in the production yield and it also becomes possible to shorten the channel and lower the threshold voltage of the MOS transistors of the memory cells. Thus, it is possible to reduce the size of memory cells and hence the overall size of the chip and improve the performance as a result of achievement of high-speed operation.

## Claims

1. A semiconductor memory circuit comprising a plurality of dynamic type memory cells which necessitate restoring of read-out data for continuing

storage of data, means for receiving an external control signal having an active level and an inactive level, means for operatively reading data from a selected memory cell in response to the active level of said external control signal and means for operatively restoring the data read from said selected memory cell into said selected memory cell in response to the inactive level of said external control signal.

2. The memory circuit according to claim 1, further comprising a plurality of word lines coupled to said memory cells, means for selecting one of said word lines in response to the active level of said external control signal, a plurality of bit lines coupled to said memory cells, a plurality of sense amplifiers coupled to said bit lines, a first reference voltage wiring coupled to said sense amplifiers, and a second reference voltage wiring coupled to said selecting means and said reading means.

3. The memory circuit according to claim 1, in which each of said memory cells includes a storage capacitor and a transfer gate transistor.

4. The memory circuit according to claim 1, in which said reading means includes a common data line, and an output circuit coupled to said common data line.

5. The memory circuit according to claim 1, further comprising a plurality of bit lines coupled to said memory cells, each of said bit lines being divided into first and second segments, and a plurality of transfer gates inserted between the first and second segments of said bit lines.

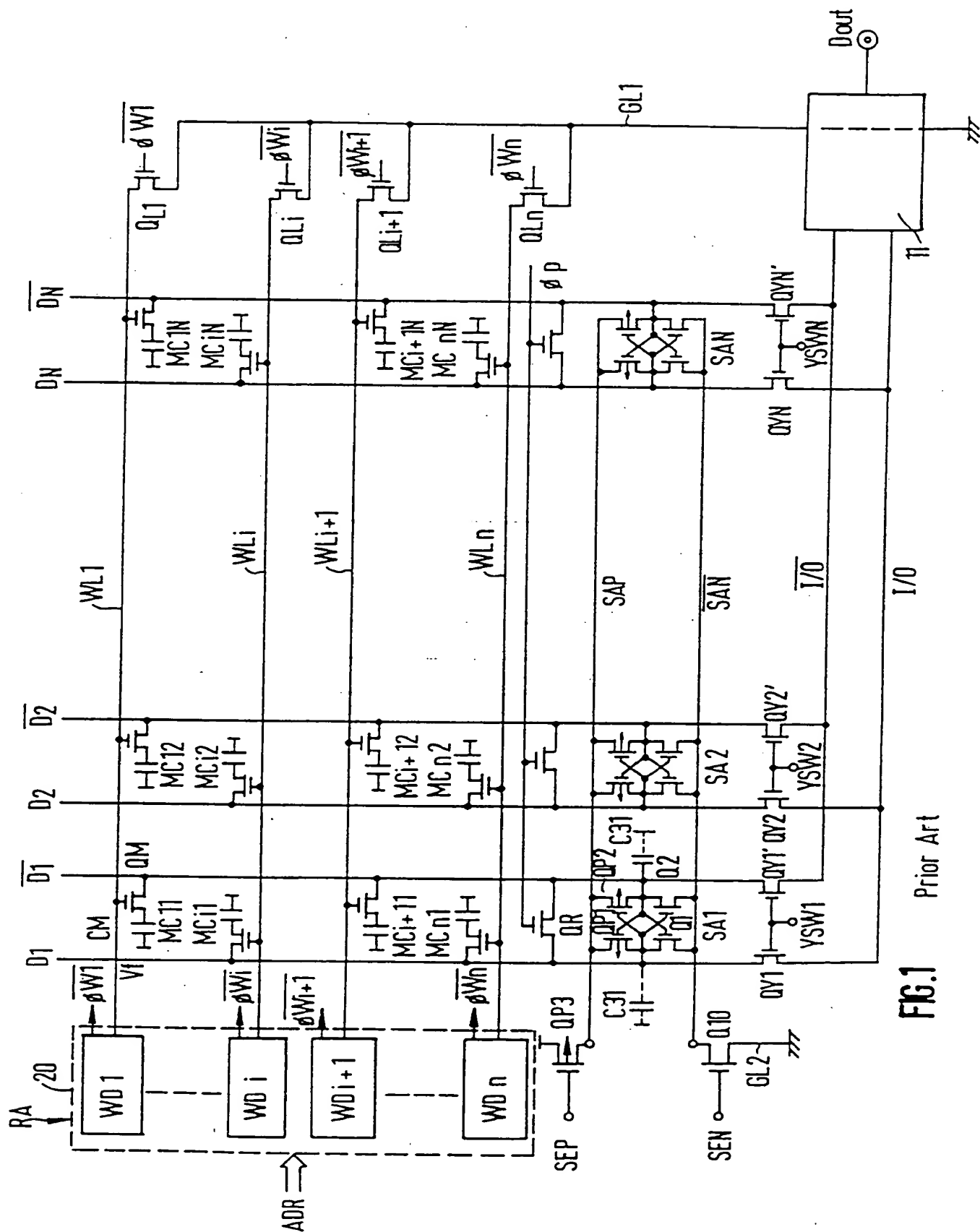


FIG. 1

Prior Art

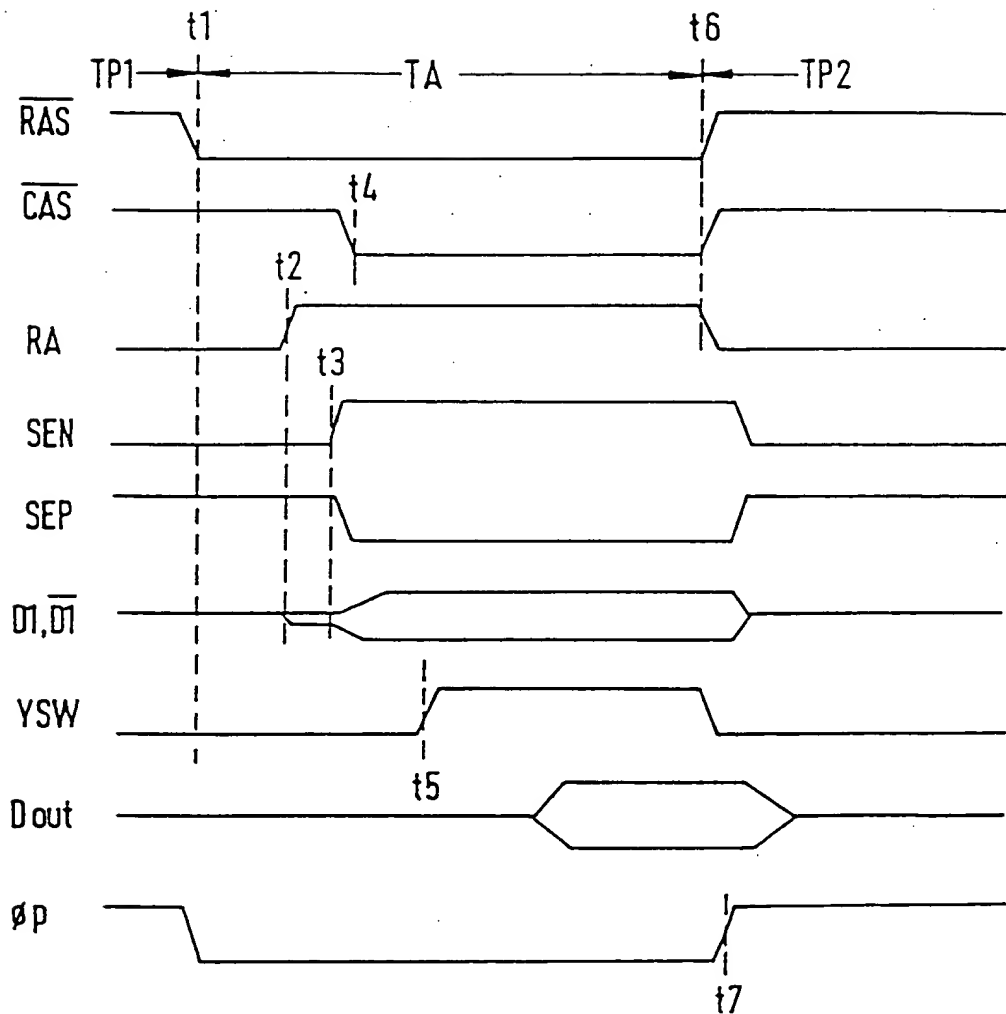


FIG.2 Prior Art



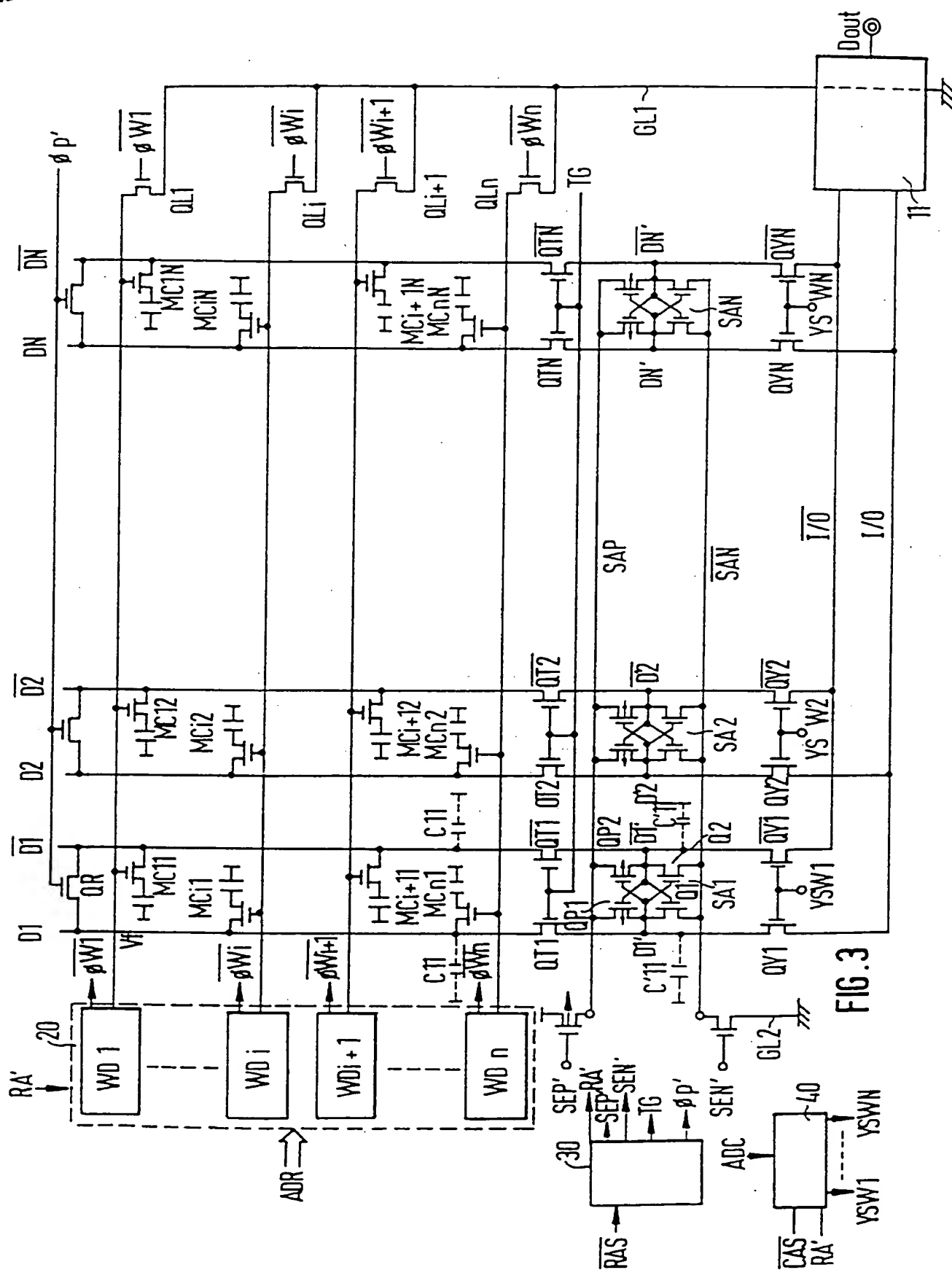


FIG. 3

1.30 0000 0000 0000  
 1.30 0000 0000 0000  
 1.30 0000 0000 0000

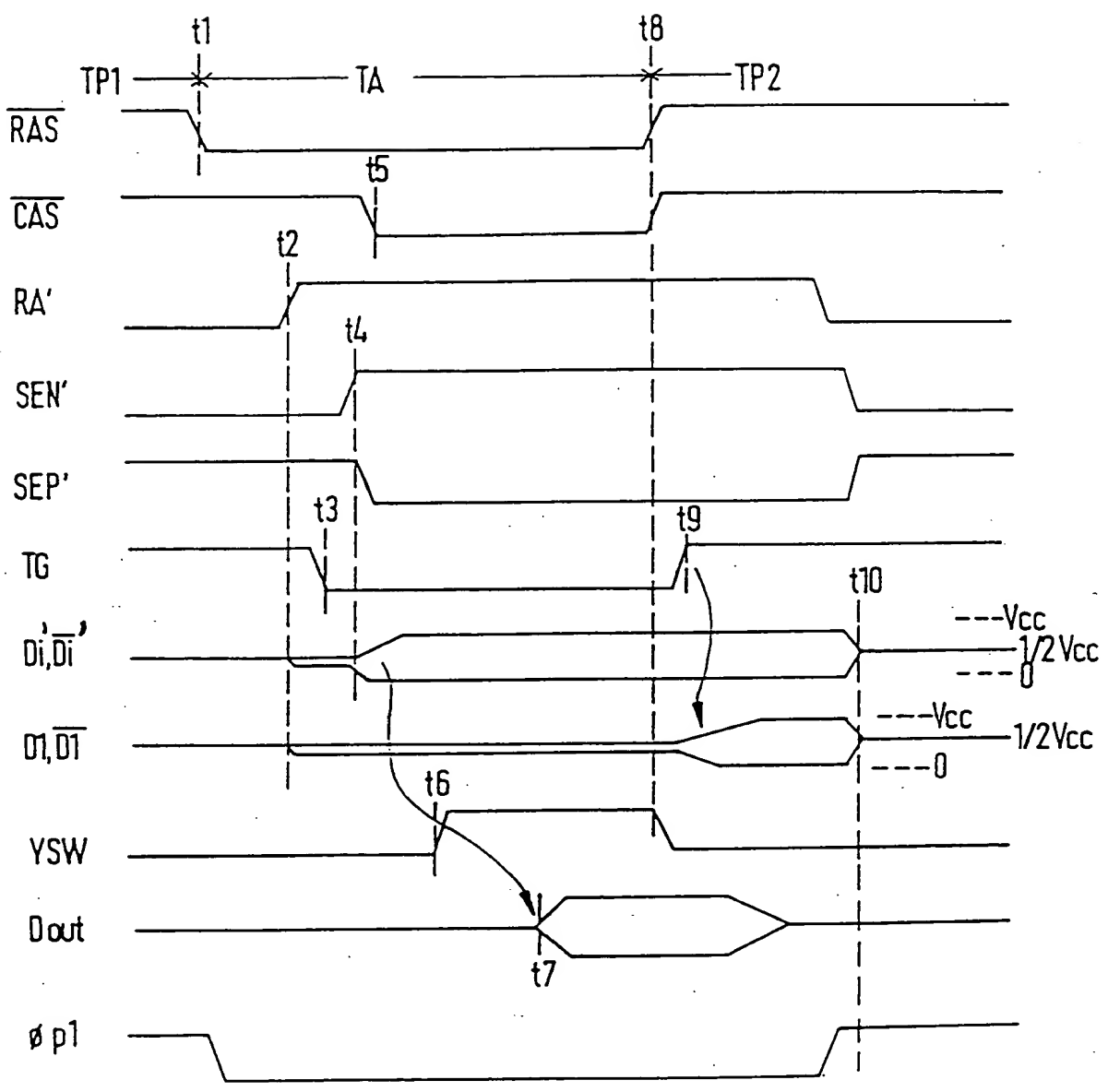
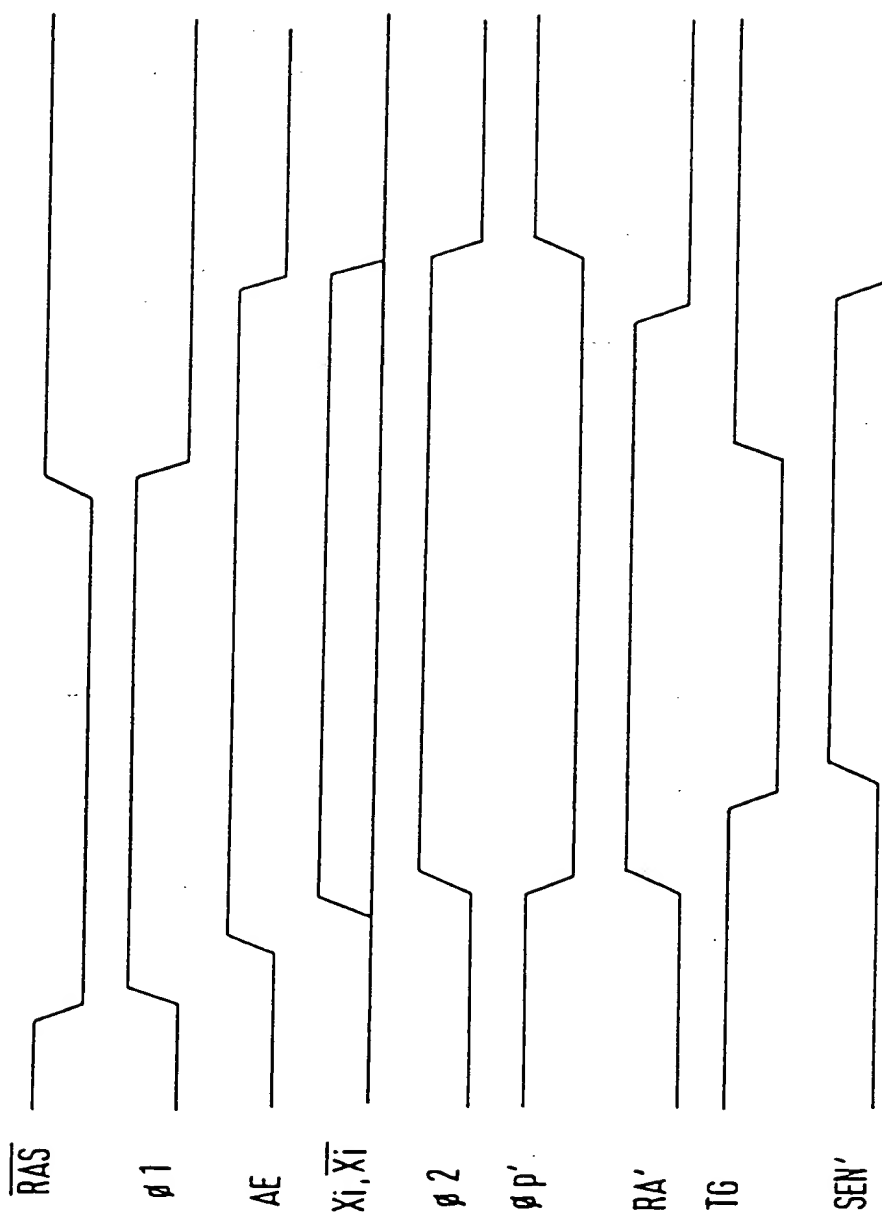


FIG. 4

The diagram illustrates a control circuit for a motor. It begins with an input  $\phi 2$  connected to an AND gate 54. The output of gate 54 is connected to an OR gate 55. The output of gate 55 is connected to a NOT gate IV4, which is then connected to an AND gate 56. The output of gate 56 is connected to a flip-flop D6. The output of flip-flop D6 is connected to a NOT gate IV6, which is then connected to an AND gate 57. The output of gate 57 is connected to a flip-flop D5. The output of flip-flop D5 is connected to a NOT gate IV8, which is then connected to an AND gate 58. The output of gate 58 is connected to a NOT gate IV9. The output of gate IV9 is connected to a transistor QP5. The output of transistor QP5 is connected to a transistor Q51. The output of transistor Q51 is connected to a transistor Q52. The output of transistor Q52 is connected to a motor (represented by a circle with a diagonal line). The circuit also includes a feedback loop from the motor output back to the input of gate 54. Other components include a NOT gate IV7, an AND gate 56, a flip-flop D6, a NOT gate IV6, an AND gate 57, a flip-flop D5, a NOT gate IV8, an AND gate 58, a NOT gate IV9, a transistor QP5, a transistor Q51, and a transistor Q52. The circuit is powered by a supply voltage  $V_{CC}$  and ground.

*[Faint, illegible handwritten notes]*



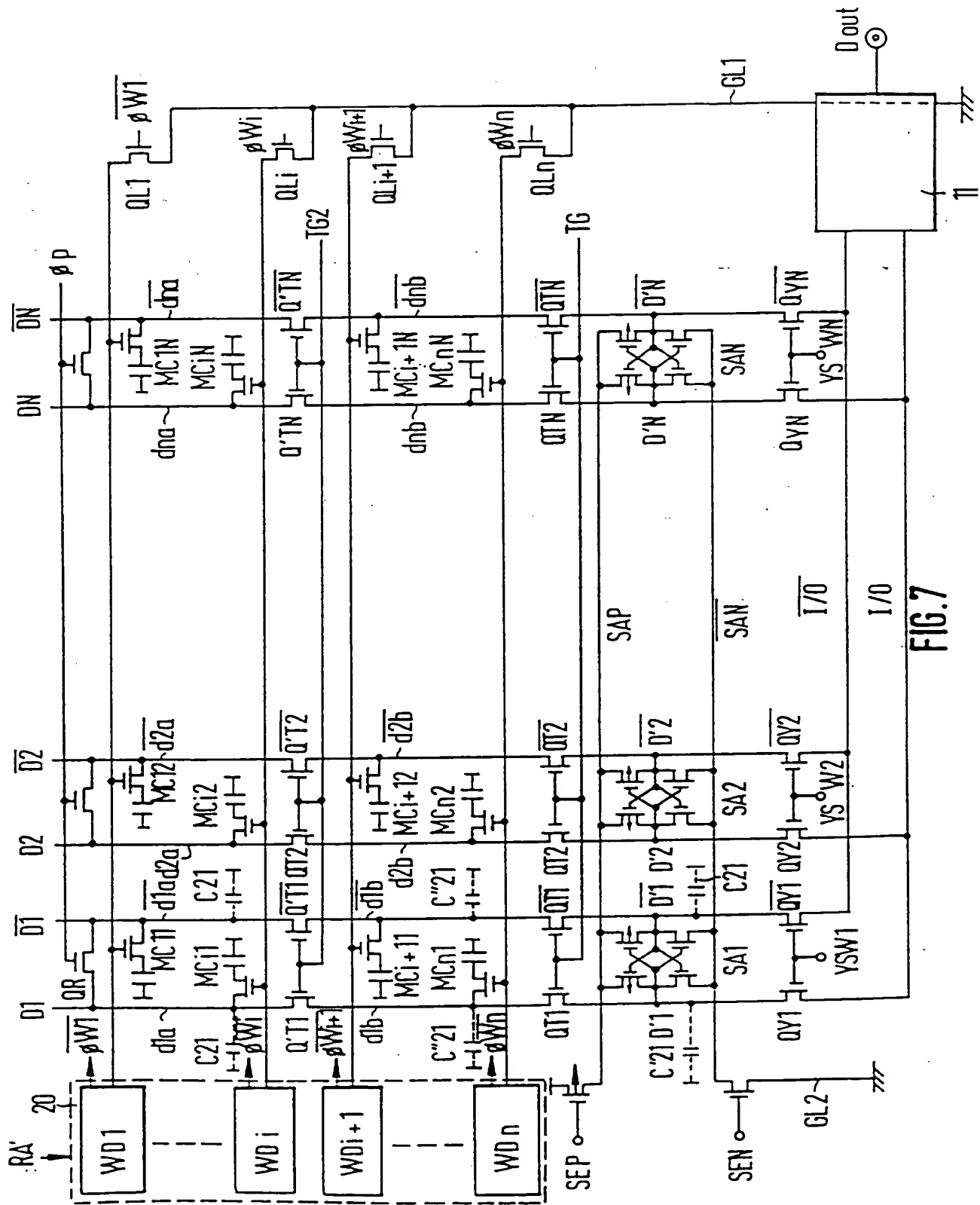


FIG. 7

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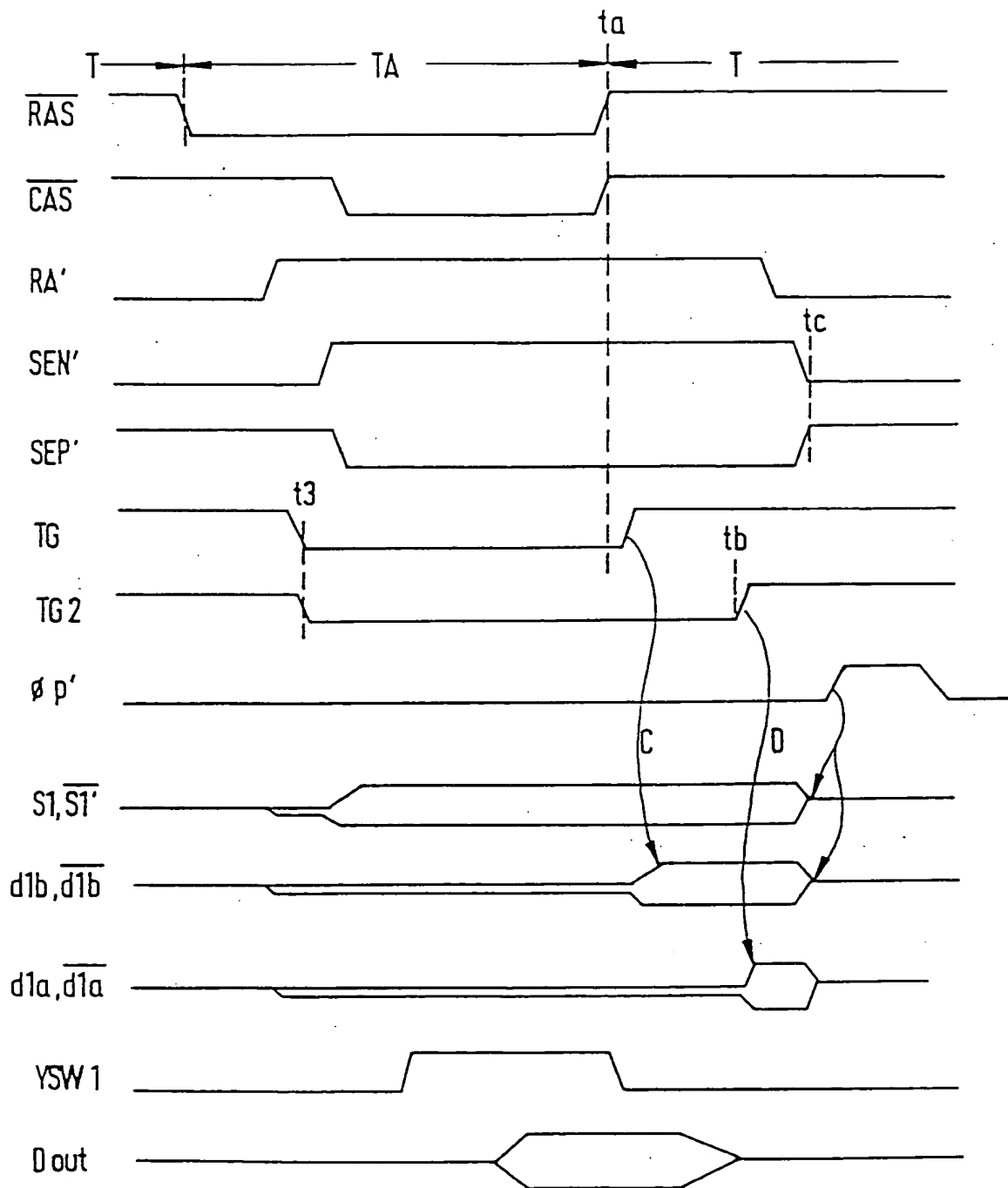


FIG. 8



(12)

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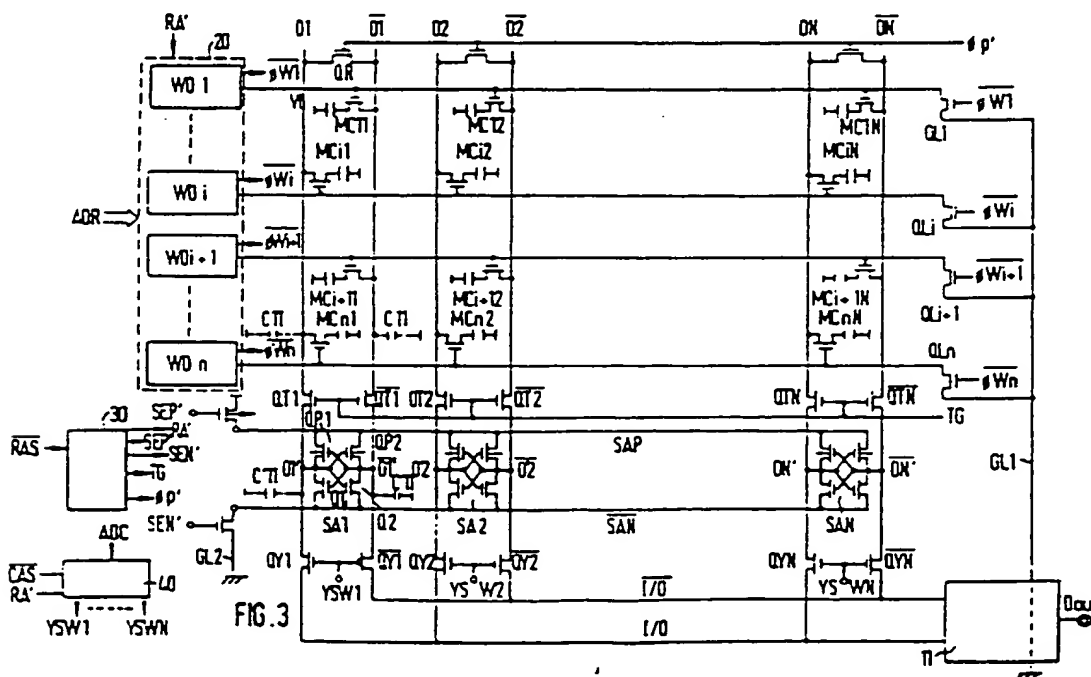
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(54) Semiconductor memory circuit having an improved restoring control circuit.

(57) A dynamic memory circuit which can operate with a reduced amount of current noise and without destruction of stored data is disclosed. The memory circuit includes dynamic memory cells necessitating restoring operation, a read circuit for performing a

read-out operation in an active state and a restore circuit for performing a restore operation in a reset state following the active state.





European  
Patent Office

## EUROPEAN SEARCH REPORT

Application Number

EP 89 12 3262

### DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int. Cl.5)
Y	EP-A-0 037 252 (FUJITSU) * Page 3, lines 1-25; figure 2 * -----	1	G 11 C 11/409
Y	EP-A-0 240 155 (FUJITSU) * Figure 2: page 8, line 26 - page 10, line 29 * -----	1	
A		3,4	
Y	WO-A-8 607 183 (AMERICAN TELEPHONE & TELE- GRAPH CO.) * Page 9, line 20 - page 10, line 27; figure 3 * -----	1	
A		3	
A	EP-A-0 185 572 (FUJITSU) * Page 4, line 25 - page 6, line 23; figures 2.3b * -----	4,5	
A	US-A-3 810 124 (W.K. JOFFMAN et al.) * Column 3, lines 9-47; figure 2 * -----	2	TECHNICAL FIELDS SEARCHED (Int. Cl.5)  G 11 C 11
P,X	EP-A-0 316 902 (NEC CORP.) * Claims 1-8 * -----	1,3-5	
The present search report has been drawn up for all claims			
Place of search  The Hague		Date of completion of search  31 January 91	Examiner  KAMSAETER K.M.S.
<div><div><b>CATEGORY OF CITED DOCUMENTS</b> X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : intermediate document T : theory or principle underlying the invention</div><div>E : earlier patent document, but published on, or after the filing date O : document cited in the application L : document cited for other reasons ..... &amp; : member of the same patent family, corresponding document</div></div>			



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